Inferring Symbolic Automata

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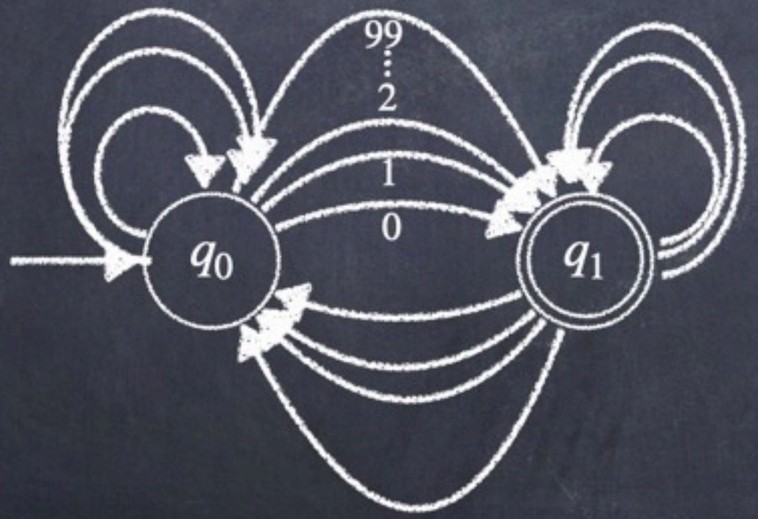
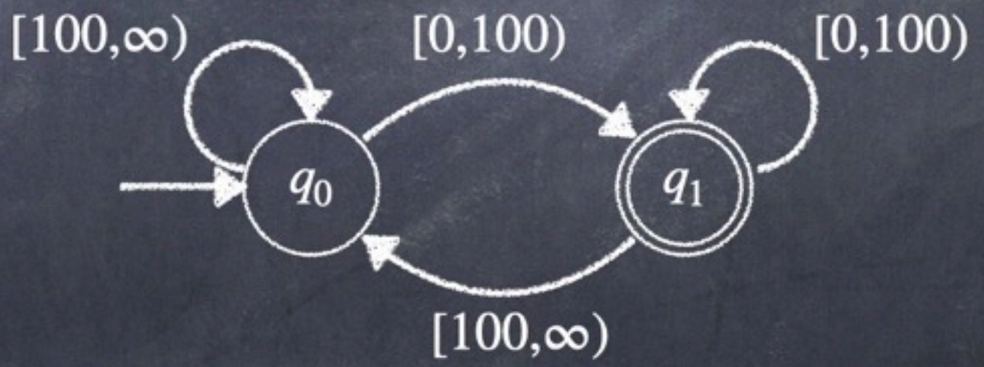
Joint work with Dana Fisman and Sandra Zilles

Symbolic Automata – SFAs

- Finite state automata
- Defined w.r.t. a Boolean algebra
- Transitions are over predicates
- Concise
- Reason about infinite domains





Monotonic Algebras

- Predicates correspond to a total order over the domain elements
- $[\psi] = \{d \mid a \leq d \leq b\}$
- Monotonic: Interval algebras (over $\mathbb{N}, \mathbb{Z}, \mathbb{R}$)

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Propositional algebra

- Predicates are Boolean combinations of atomic propositions
- $(p_1 \wedge p_2) \vee (\neg p_1 \wedge p_3)$
- Not monotonic!

Automata Learning

- Active learning - L* style learning [Angluin 1987]

- Passive learning

Automata Learning

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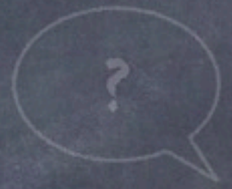
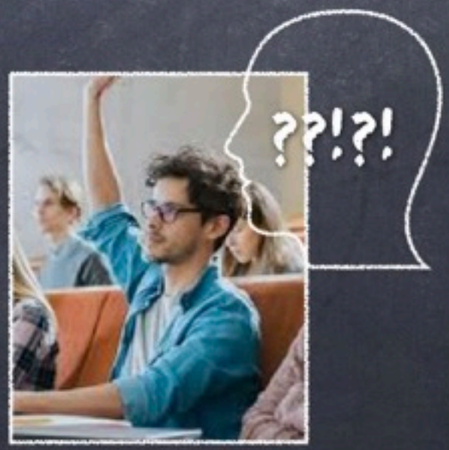


- Passive learning

$\langle w_1, \bot \rangle$
$\langle w_2, \top \rangle$
$\vdots$
$\langle w_n, \bot \rangle$



Learnability of a Class of Languages via Representation $\mathscr{R}$

- Different representations of languages
  - E.g. regular languages – DFAs, NFAs
- A class of languages $\mathscr{L}$ is learnable via representation in $\mathscr{R}$ if there is an algorithm ALG such that we can apply ALG to learn a representation in $\mathscr{R}$ for every language in $\mathscr{L}$
  - E.g. L* algorithm for regular languages via representation in DFAs

Active Learning of SFAs

- Positive results
  - Learning of SFAs over monotonic algebras using membership and equivalence queries
    - [Maler & Mens 2014], [Maler & Mens 2017]
    - [Chubachi, Diptarama, Yoshinaka, Shinohara 2017]
  - MAT* algorithm for learning SFAs
    - [Argyros & D'Antoni 2018]

Active Learning of SFAs

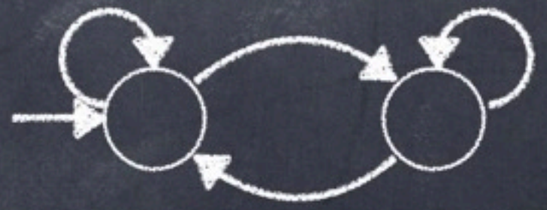
- First negative result
- Necessary condition:
- We can polynomially learn SFAs over a Boolean algebra $\mathcal{A}$ using membership and equivalence queries only if we can polynomially learn the predicates of $\mathcal{A}$ using membership and equivalence queries

→ SFAs over the propositional algebra are not polynomially learnable

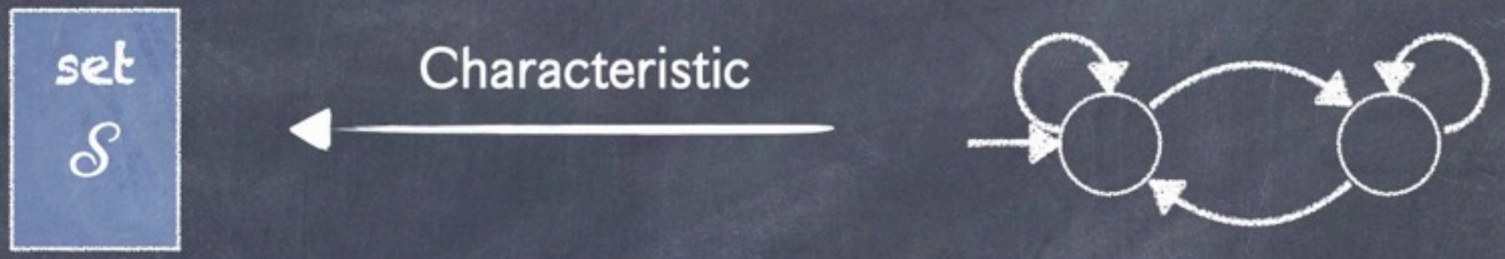
Passive Learning

set
S'

Infer



Identification in the Limit Using Polynomial Time and Data



[Gold 1978, de la Higuera 1997]

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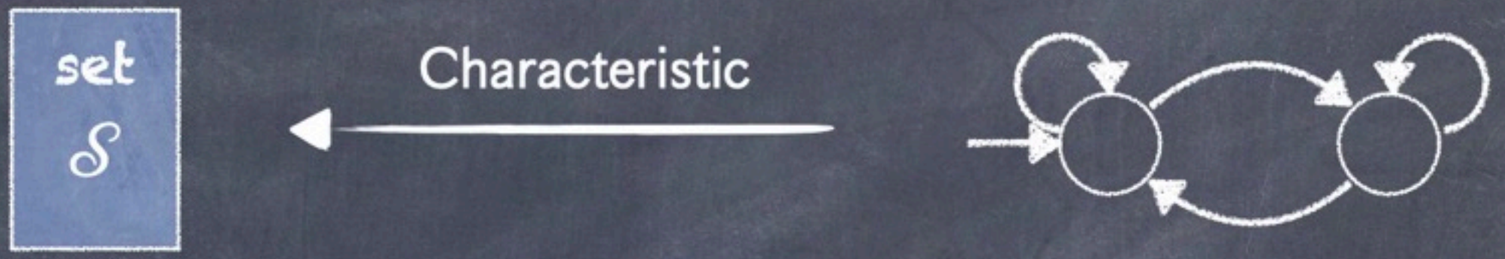

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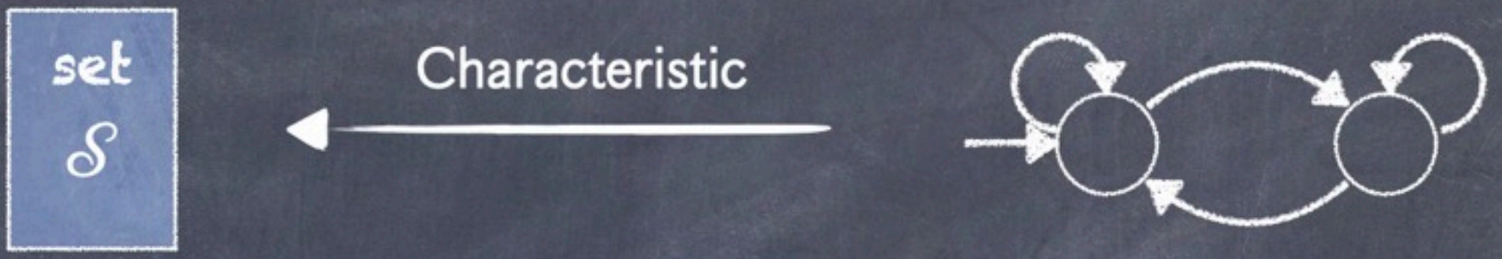

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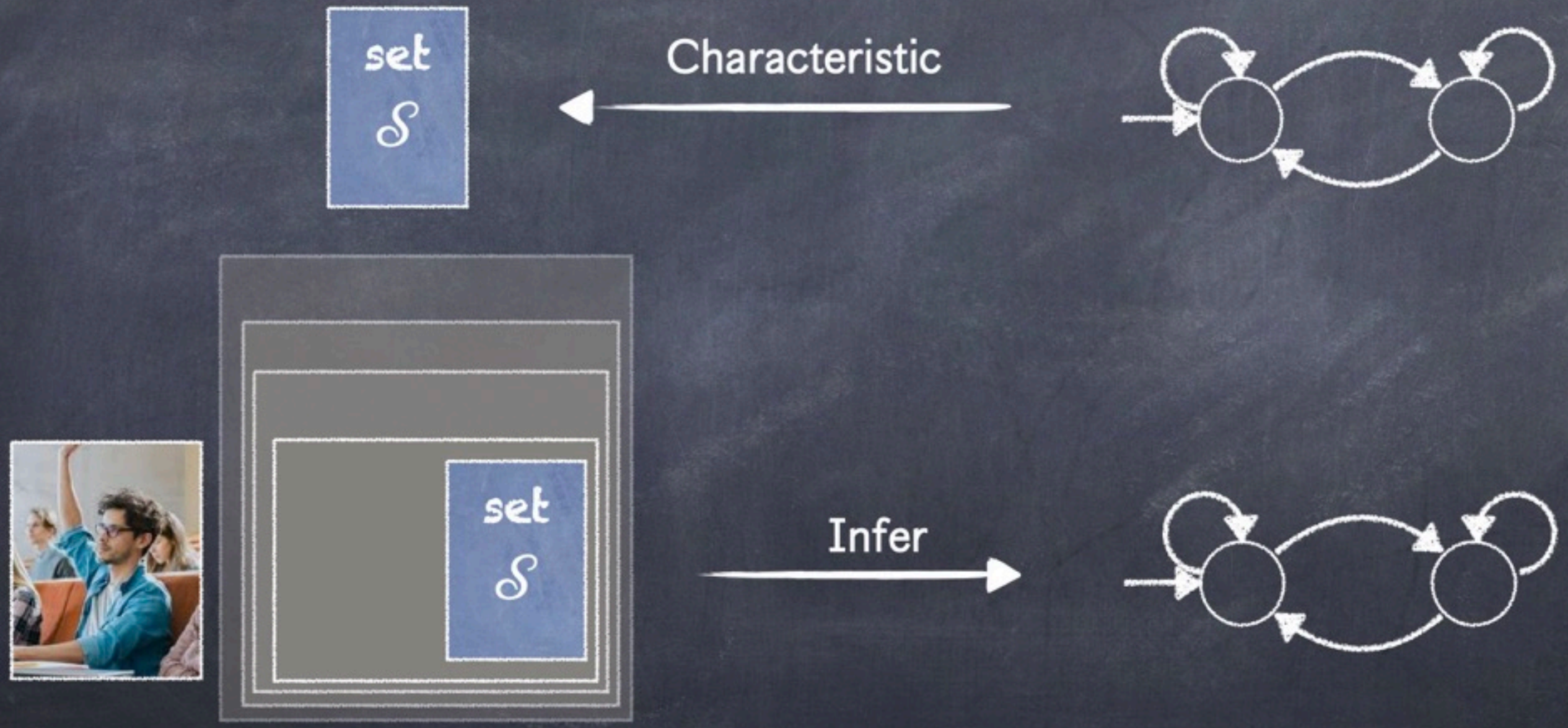

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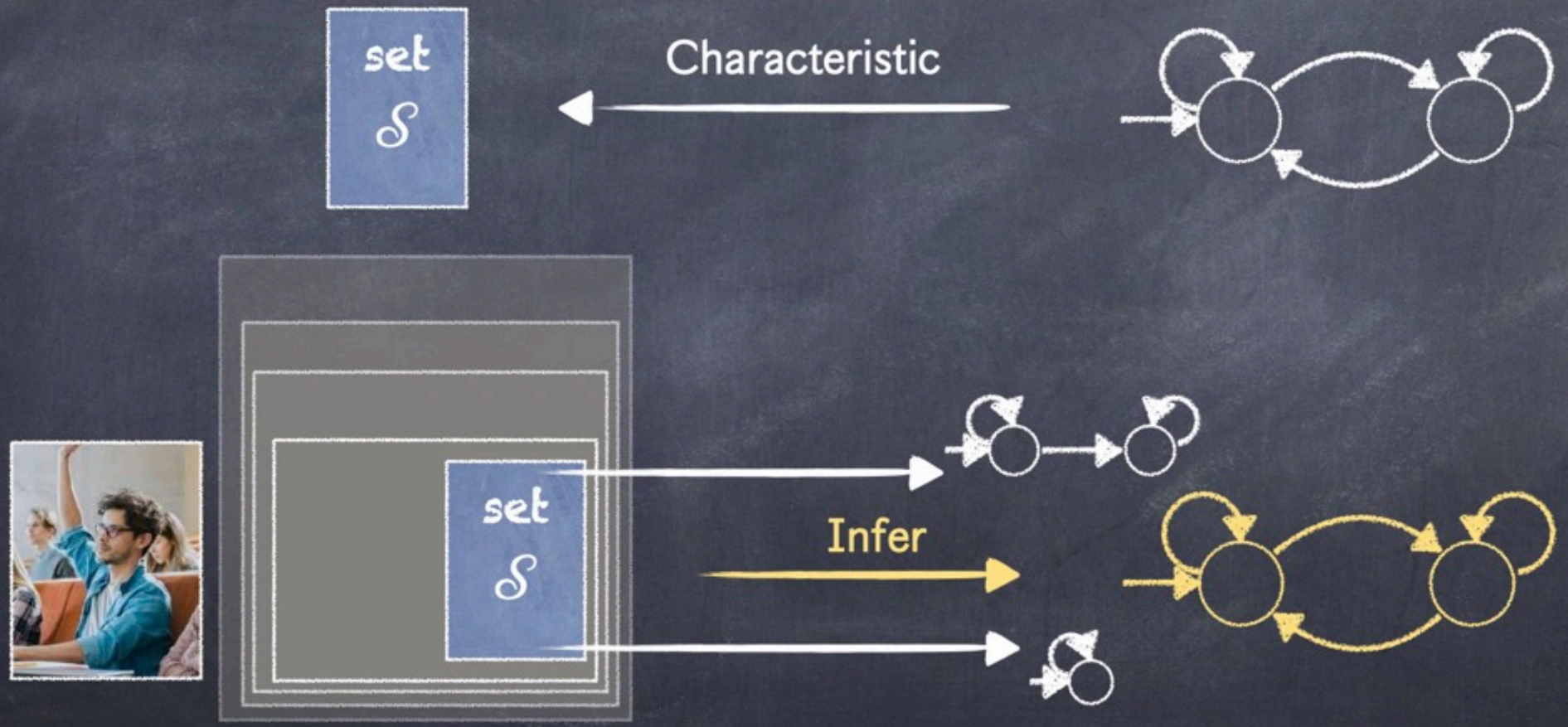

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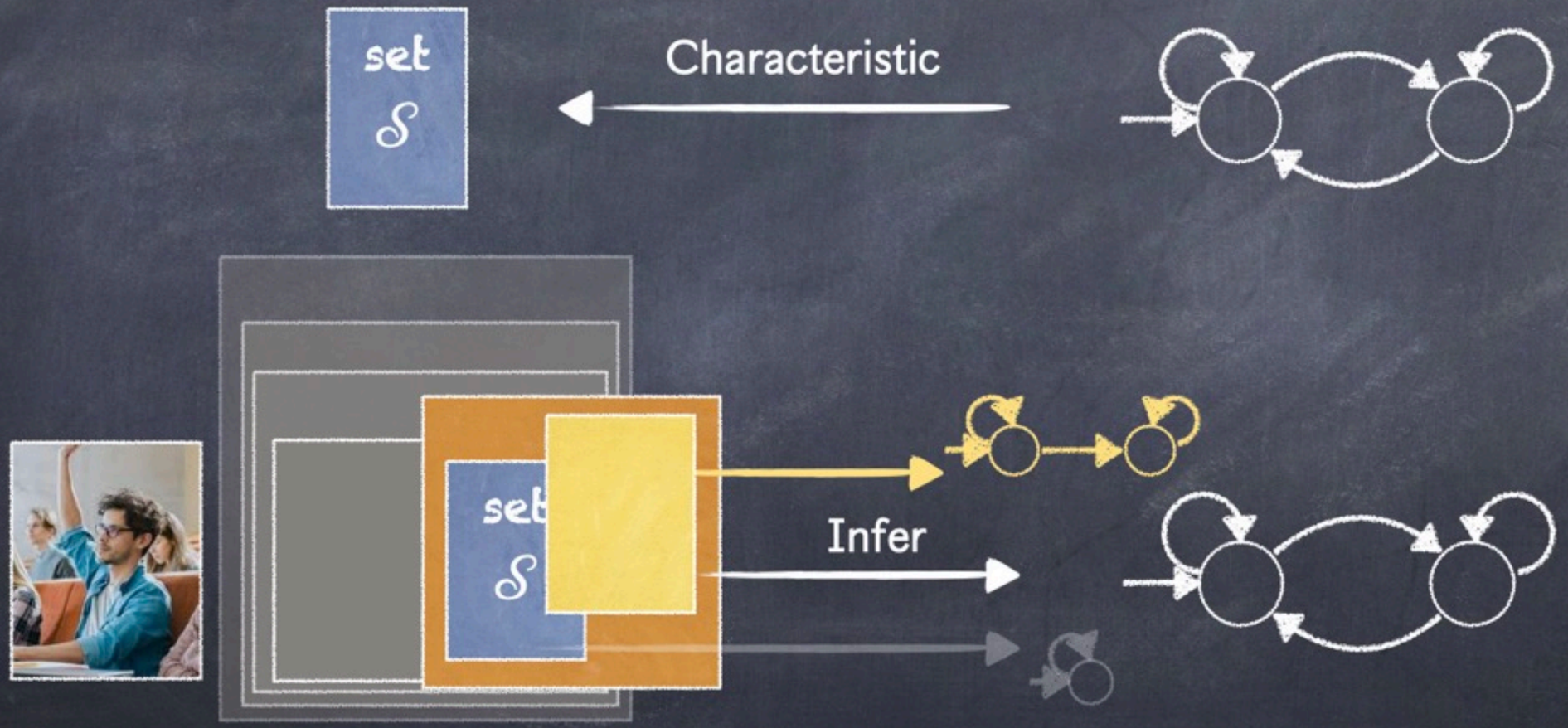

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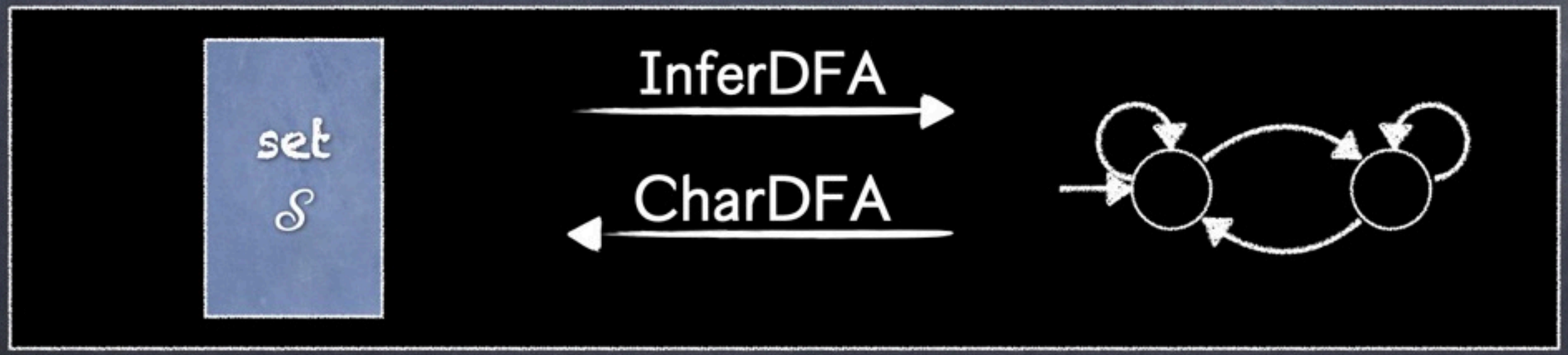
Identification in the Limit for DFAs



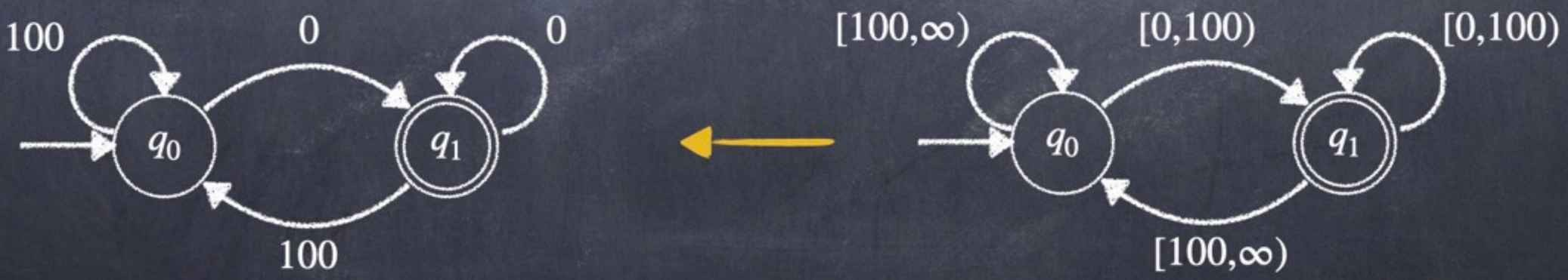
Concrete Sample set

DFA

[Oncina&García 1992]

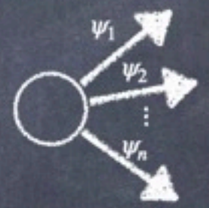
Identification in the Limit for SFAs - CharSFA

- Learning with respect to the concrete alphabet
- Creating a set of concrete words

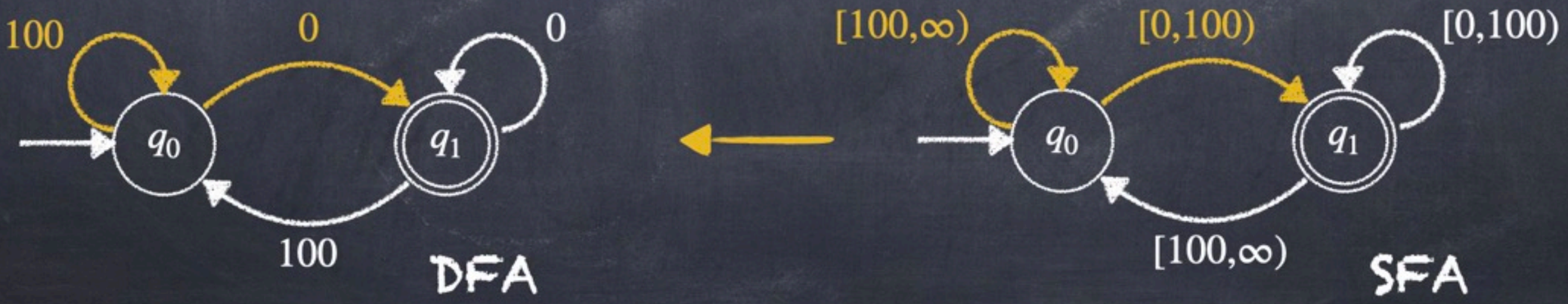


Identification in the Limit for SFAs - CharSFA

- Learning with respect to the concrete alphabet
- Creating a set of concrete words
- concretize($\langle \psi_1, \ldots, \psi_n \rangle$) = $\langle \Gamma_1, \ldots, \Gamma_n \rangle$



- concretize($\langle [0,100), [100,\infty) \rangle$) = $\langle \{0\}, \{100\} \rangle$



Identification in the Limit for SFAs - CharSFA

- Creating a set of concrete words

Identification in the Limit for SFAs - InferSFA

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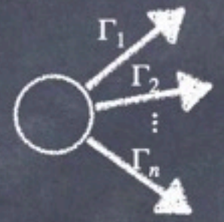
Concrete Sample set

DFA Generalize

Identification in the Limit for SFAs - InferSFA

- generalize($\langle \Gamma_1, \ldots, \Gamma_n \rangle$) = $\langle \psi_1, \ldots \psi_n \rangle$ 

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Identification in the Limit for SFAs - InferSFA

Concrete Sample set

DFA Generalize

Identification in the Limit for SFAs - InferSFA

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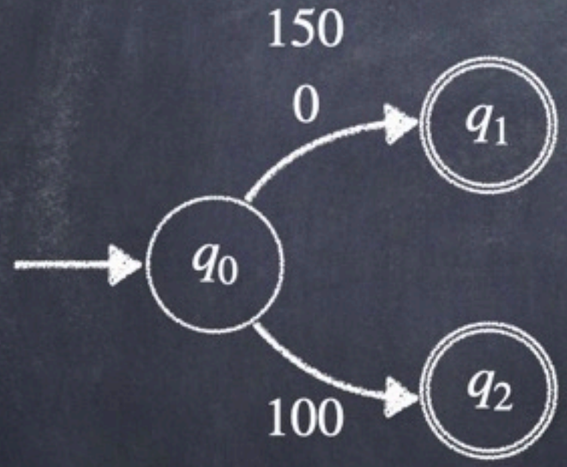
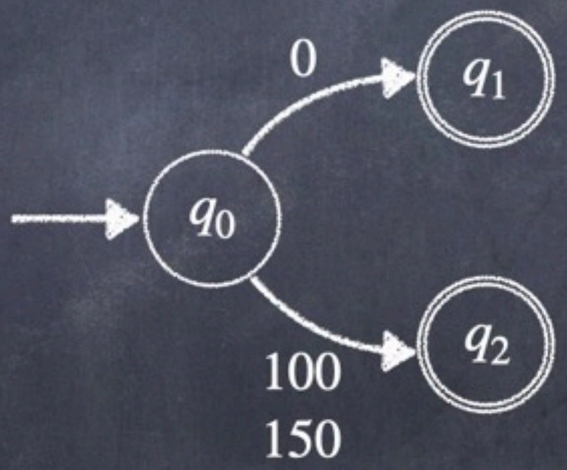
Concrete Sample set

DFA

Identification in the Limit for SFAs - InferSFA

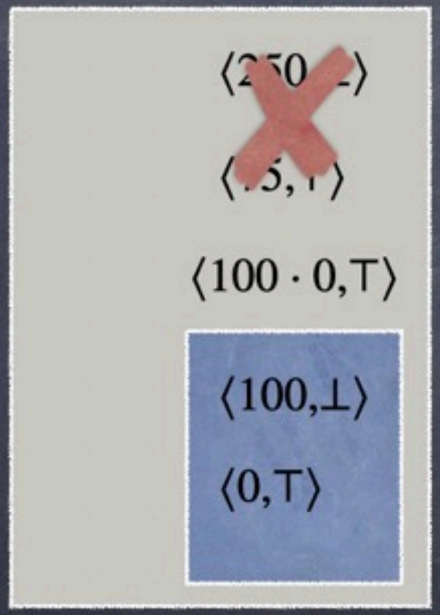
- Additional alphabet adds confusion



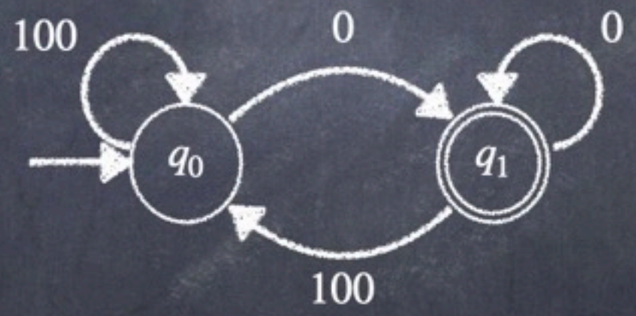


$\langle \epsilon, \perp \rangle, \langle 0, \top \rangle, \langle 100, \top \rangle$
$\langle 0 \cdot 0, \top \rangle, \langle 100 \cdot 0, \perp \rangle$
$\langle 150, \top \rangle$

Identification in the Limit for SFAs - InferSFA



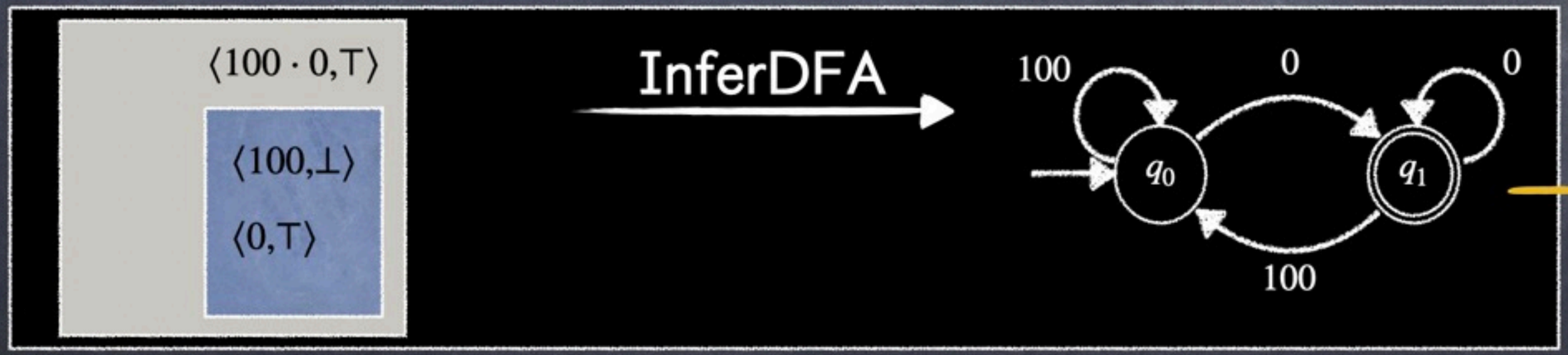
Decontaminate



Concrete Sample set

DFA

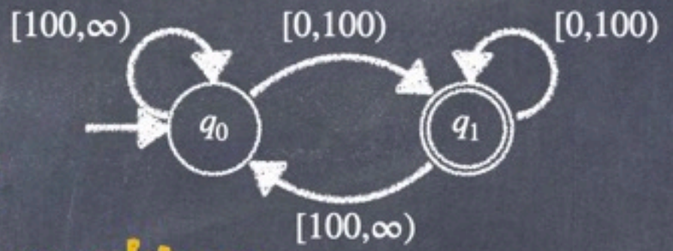
Identification in the Limit for SFAs - InferSFA



Concrete Sample set

DFA Generalize

The Whole Process



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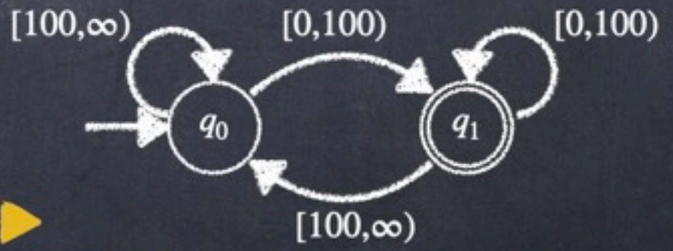
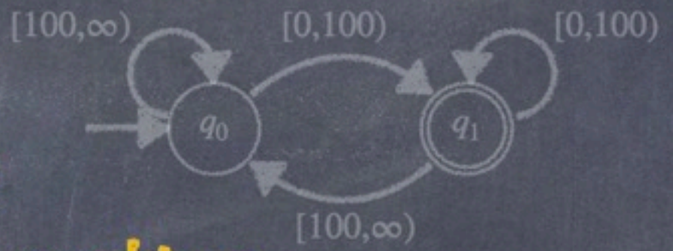


Generalize

Sufficient Condition

Monotonic algebras



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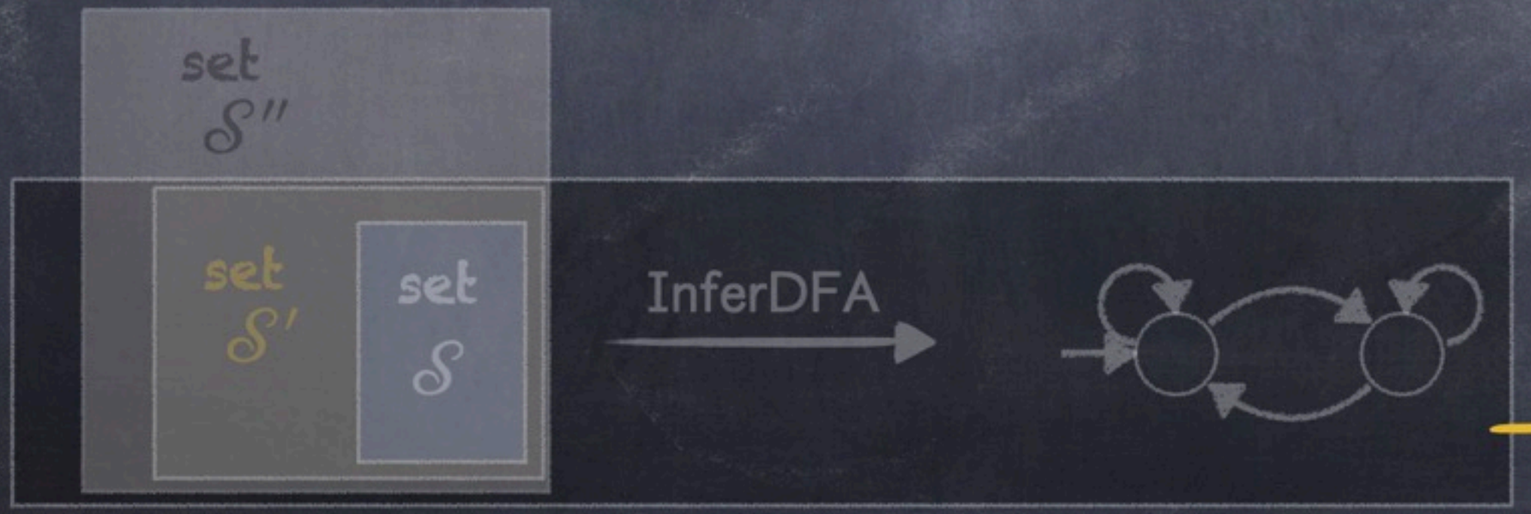
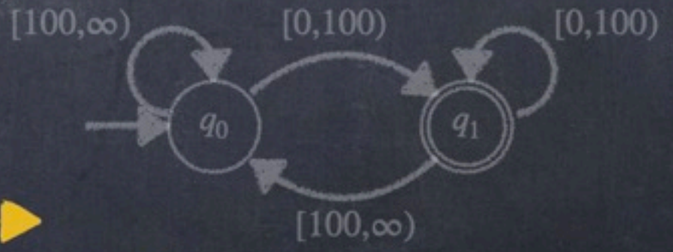
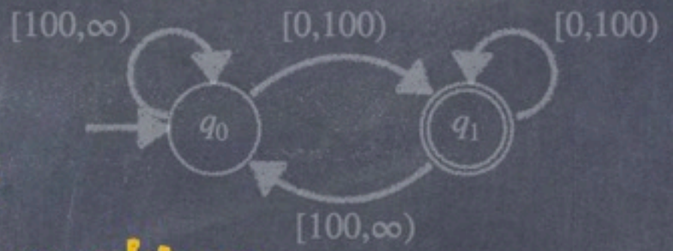




Generalize

Necessary Condition

Propositional Algebra



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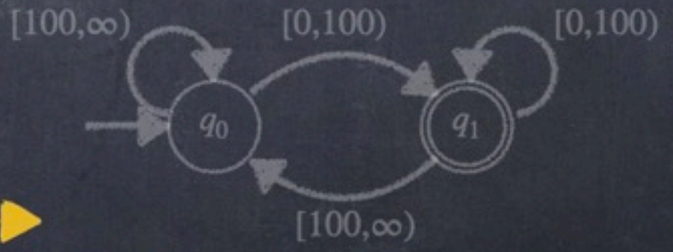

Generalize

Summary

- Active Learning
  - Necessary condition
  - SFAs over the propositional algebra are not polynomially learnable

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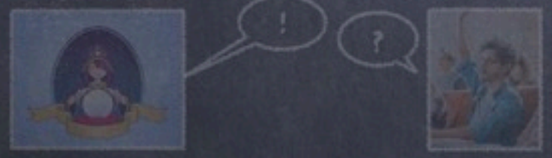
  - Necessary condition
  - SFAs over the propositional algebra are not polynomially learnable

- Passive Learning

  - Necessary condition & sufficient condition for learning SFAs
  - SFAs over the propositional algebra are not polynomially learnable
  - Learning algorithm for SFAs over monotonic algebras
  - Learning scheme for the paradigm of identification in the limit of SFAs

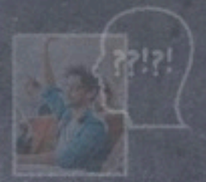
THANK YOU!

- Active Learning
  - Necessary condition
  - SFAs over the propositional algebra are not polynomially learnable



Questions?

- Passive Learning



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